#### Dynamic Position Auctions with Consumer Search

#### Scott Duke Kominers

Harvard University

#### Algorithmic Aspects in Information and Management June 16, 2009

Background

### What are position auctions?

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Background

What are position auctions?

#### Definition

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• Position auctions are used to allocate sponsored search links to advertisers!

Background

### Why study position auctions?

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- Sponsored search is a multibillion-dollar industry
- The mechanisms used are relatively new
- Welfare implications not well-understood

Background

#### **Previous Position Auction Models**

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Exogenous Click-through Rates

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Our Model

#### Framework & Conventions

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Our Model

### Framework & Conventions

#### Athey and Ellison (2008) Model

Our Model

### Framework & Conventions

#### Athey and Ellison (2008) Model • *N* advertisers

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#### Our Dynamic Model

• Extends Athey and Ellison (2008)

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  - Sequential rounds
  - Synchronous updating
  - Advertisers play a "best-response" strategy
  - Consumers ignorant of dynamics

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- Unique fixed point
  - Athey and Ellison (2008) Envy-Free Equilibrium

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### Main Result

Image: A matrix and a matrix

# Main Result Theorem (Convergence Theorem)

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If all advertisers play the Restricted Balanced Bidding strategy, then their bids converge to the fixed point; this convergence is efficient.

• The dynamic model is "well-approximated" by the static model.

#### Parameters

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Our Model

### Results

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### Lemma

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• Within  $t_1$  rounds, the N - M lowest-quality advertisers "drop out" of contention.

### Convergence of the *M* Positions • By the Lemma

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  - Set of advertisers in positions of P:  $\pi(P)$
  - Next round, all advertisers in  $\pi(P)$  repeat their bids.
    - If  $\pi(P) = \{1, \dots, M\}$ , then we are done.

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- Set of stable positions:  $P = \{p + 1, \dots, M\}$
- Advertiser  $\pi \notin \pi(P)$  with the lowest bid
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- (Significant divergence from Cary et al. (2008))

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Our Model

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#### Lemma

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$$\epsilon = \frac{G(\bar{q}_M)}{2G(\bar{q}_1)} (1 - \gamma^{**}) \min_{\phi \neq \phi'} |q_\phi - q_{\phi'}| \left( \prod_{j=1}^M (1 - q_j) \right).$$

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At most  $\log_{1/\gamma^{**}}((q_1 - q_{M+1})/\epsilon)$  consecutive instances of  
Case 3 may occur.

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 $\bullet \ \Rightarrow \mathbb{QED}$ 

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• This also yields probability-1 efficient convergence in an asynchronous bidding model.

Discussion

#### **Possible Generalizations**

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#### • Our method pprox Cary et al. (2008)'s method

 Our method ≈ Cary et al. (2008)'s method; its applicability is naïvely surprising.

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Three key steps:

Three key conditions:

• Our method pprox Cary et al. (2008)'s method

Three key steps:

restriction of the strategy space

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- analysis of low-quality advertisers' behaviors

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Three key steps:

- restriction of the strategy space
- analysis of low-quality advertisers' behaviors
- **o** proof that the *M* positions stabilize
- Three key conditions:
  - unique envy-free equilibrium
  - Iow-quality advertisers drop out efficiently
  - monotone equilibrium strategy

Discussion

## Conclusion

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Image: A mathematical states of the state

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## Conclusion

Convergence should be demonstrable in dynamic position auction models with sufficiently well-behaved static equilibrium strategies.

# Questions?

QED

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**Dynamic Position Auctions** 

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